

An overview on control theory for (max,plus) linear systems,

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Angers, along the Loire Valley



Angers, along the Loire Valley



- 1H30 in south west of Paris
- 350000 people, 16th city of France
- 1 State university, with 32000 students,



- Denis Papin graduated from University of Angers in 1669



- Carlos Andrey Maia graduated in 2004

Links between UFMG and UA, more than 10 years old



UFMG - UA History in some points

- PhD of Carlos Andrey Maia, in [2001-2004]
- CAPES - Cofecub Project [2008-2012]
- Brafitec Project [2012-2016]
- Vinicius Mariano Gonçalves PhD defense, [2013-2014] in Angers



- Fields medal 1966
- Member of the Nicolas Bourbaki Seminar

An overview on control theory for (max,plus) linear systems

Outline

- **Introductory examples**
- $(\max,+)$ algebra in few words
- Discrete event systems in $(\max,+)$ algebra
- Optimal control on a finite horizon
- Optimal filtering
- Optimal closed loop control, how to take into account disturbances?
- Sub observer synthesis
- Interval arithmetic and semiring, robust controller synthesis

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(max,+) algebra in few words

Idempotent Semiring

- Sum \oplus , associative, commutative, zero element denoted ε ,
- Product \otimes , associative, identity element denoted e ,
- Product \otimes distributes with respect of sum,
 $(a \oplus b) \otimes c = a \otimes c \oplus b \otimes c$,
- Zero element ε is absorbing, $a \otimes \varepsilon = \varepsilon$
- The sum is idempotent, $a \oplus a = a$.
- $a \oplus b = a \Leftrightarrow b \preceq a \Leftrightarrow a \wedge b = b$
hence a semiring has a complete lattice structure, with ($\varepsilon = \perp$) as bottom element and ($\top = \bigoplus_{x \in \mathcal{S}} x$) as top element.

Example : (max,+) algebra, $\overline{\mathbb{Z}}_{\max}$, Cuninghame-Green 1962

► More

Sum \oplus is the operator *max*, product \otimes is classical sum $+$, $\varepsilon = -\infty$ and $e = 0$, then :

$$\begin{aligned}1 \oplus 1 &= 1 = \max(1, 1), \\2 \otimes 1 &= 3 = 2 + 1.\end{aligned}$$

$(\max, +)$ algebra in few words

Residuation Theory (Croisot 1956, Blyth 1972)

It allows to define a kind of inverse for order preserving mapping, defined over ordered sets. The idempotent law \oplus induces that a semi ring has an order structure, *i.e.*, $a \oplus b = a \Leftrightarrow b \preceq a$, then this theory is suitable to consider inversion problem over idempotent semi ring.

Inequality $a \otimes x \preceq b$

Over a complete idempotent semi ring, inequality $a \otimes x \preceq b$ admits a greatest solution, denoted, $x = a \backslash b$.

Example : $(\max, +)$ algebra \mathbb{Z}_{\max}

Inequality $5 \otimes x \preceq 3$ admits a greatest solution $x = 5 \backslash 3 = 3 - 5 = 2$. It achieves equality in the scalar case.

$(\max, +)$ algebra in few words

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(max,+) algebra in few words

Fixed point equation $x = ax \oplus b$

► More

Theorem : Over a complete idempotent semi ring \mathcal{T} , the least solution to $x = ax \oplus b$ is $x = a^*b$ with $a^* = \bigoplus_{i \geq 0} a^i = e \oplus a \oplus a^2 \oplus \dots$

* is called Kleene star operator.

We will denote \mathcal{M} the mapping defined over \mathcal{T} s.t. $\mathcal{M} : x \mapsto x^*$

Inequality $x^* \preceq a$

► More

Theorem : Mapping \mathcal{M} restricted to its image is residuated.

Application : If $a \in \text{Im}\mathcal{M}$ then inequality $x^* \preceq a$ admits a greatest solution $x = a^*$

(max,+) algebra in few words

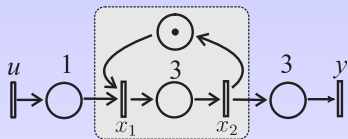
Matrix

Let A, B, C three matrices in $\overline{\mathbb{Z}}_{\max}^{n \times n}$

- $(A \oplus B)_{ij} = A_{ij} \oplus B_{ij}$
- $(A \otimes B)_{ik} = \bigoplus_{j=1 \dots n} (A_{ij} \otimes B_{jk})$
- $(A \oslash B)_{ik} = \bigwedge_{j=1 \dots n} (A_{ji} \oslash B_{jk})$, where $A \oslash B$ is the greatest such $AX \preceq B$

▶ More

TEG Model in $\overline{\mathbb{Z}}_{\max}$



Firing Date [Cohen et al., 85]

$x_i(k)$: date of the firing numbered k
for the transition labelled i .

For each transition :

$$x_1(k) = \max(1 + u(k), x_2(k - 1))$$

$$x_2(k) = 3 + x_1(k)$$

$$y(k) = 3 + x_2(k)$$

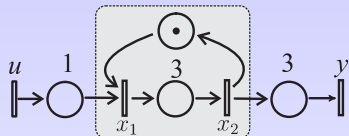
In $\overline{\mathbb{Z}}_{\max}$

► $\overline{\mathbb{Z}}_{\max}$:

$$x_1(k) = 1 \otimes u(k) \oplus x_2(k - 1)$$

$$x_2(k) = 3 \otimes x_1(k)$$

$$y(k) = 3 \otimes x_2(k)$$



Firing Date [Cohen et al., 85]

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Dynamic Model

$$x(k) = Ax(k-1) \oplus Bu(k)$$

$$y(k) = Cx(k)$$

TEG Model in $\overline{\mathbb{Z}}_{\max}[\gamma]$

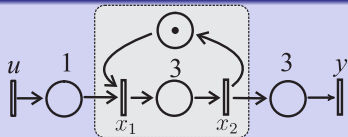
γ transform [Cohen et al,89]

► More

γ transform of $x(k)$ is a formal series $x(\gamma) = \bigoplus_{k \in \mathbb{Z}} \gamma^k x(k)$.

γ -operator is like a backward shift operator in the event domain,
 $y(\gamma) = \gamma x(\gamma) \Leftrightarrow y(k) = x(k-1) \forall k$.

The previous system in $\overline{\mathbb{Z}}_{\max}[\gamma]$:



$$x_1(\gamma) = 1u(\gamma) \oplus \gamma x_2(\gamma)$$

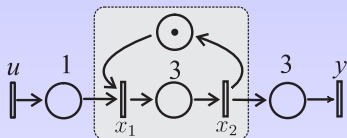
$$x_2(\gamma) = 3x_1(\gamma)$$

$$y(\gamma) = 3x_2(\gamma)$$

General case , transfer matrix

$$\begin{cases} x(\gamma) = Ax(\gamma) \oplus Bu(\gamma) \\ y(\gamma) = Cx(\gamma) \end{cases} \Rightarrow y = CA^*Bu = Hu$$

TEG Model in $\overline{\mathbb{Z}}_{\max}[\gamma]$



The previous system in $\overline{\mathbb{Z}}_{\max}[\gamma]$:

$$\begin{aligned}
 x(\gamma) &= Ax(\gamma) \oplus Bu(\gamma) = \begin{pmatrix} \varepsilon & \gamma \\ 3 & \varepsilon \end{pmatrix} x(\gamma) \oplus \begin{pmatrix} 1 \\ \varepsilon \end{pmatrix} u(\gamma) \\
 y(\gamma) &= Cx(\gamma) = \begin{pmatrix} \varepsilon & 3 \end{pmatrix} x(\gamma)
 \end{aligned}$$

The previous system in $\overline{\mathbb{Z}}_{\max}[\gamma]$:

► More

$$\begin{aligned}
 x(\gamma) &= A^*B = \begin{pmatrix} (3\gamma)^* & \gamma(3\gamma)^* \\ 3(3\gamma)^* & (3\gamma)^* \end{pmatrix} \begin{pmatrix} 1 \\ \varepsilon \end{pmatrix} u(\gamma) \\
 y(\gamma) &= CA^*B = (7(3\gamma)^*) u(\gamma)
 \end{aligned}$$

Operations over $\overline{\mathbb{Z}}_{\max}[\gamma]$

Operations over $\overline{\mathbb{Z}}_{\max}[\gamma]$

- $s = s_1 \oplus s_2$ ▶ More
- $s = s_1 \otimes s_2$ ▶ More
- $s = s_1 \wedge s_2$, asymptotic slope $\sigma_{\infty}(s) = \max(\sigma_{\infty}(s_1), \sigma_{\infty}(s_2))$
- s^* , asymptotic slope $\sigma_{\infty}(s^*) = \min(\sigma_{\infty}((p_i)), \sigma_{\infty}(q_i), \sigma_{\infty}(r))$
- $s = s_1 \setminus s_2$, if $\sigma_{\infty}(s_2) \leq \sigma_{\infty}(s_1)$ then $\sigma_{\infty}(s) = \sigma_{\infty}(s_2)$, else $s = \varepsilon$

Problem Formulation :

Let z be a desired output. Is it possible to a control input in order to get an output y as close as possible to z while respecting the constraint : $y \preceq z$.

Solution

The optimal input is given by :

$$u_{opt} = (CA^*B) \oslash z$$

It is the greatest input which achieves the inequality :

$$y = (CA^*B)u_{opt} \preceq z$$

In manufacturing setting z corresponds to the customer demands, u the input of raw parts in the system, y the output of processed parts. Optimal control u_{opt} is the one which minimizes the internal stock while ensuring the customer demand is honored.

Practical computation :

Residuation of matrices of series... or In (max,plus) algebra it can be formally expressed as :

$$\zeta(k) = A \backslash \zeta(k+1) \wedge C \backslash z(k)$$

$$u(k) = B \backslash \zeta(k)$$

It is a back tracking computation , very reminiscent to the optimal control for classical dynamical systems. $\zeta(k)$ is the adjoint state of the system.

Drawback :

Desired output has to be known a priori, or a receding horizon has to be performed.

Optimal Filtering, model matching problem : JESA, Hardouin et al. 1995

Problem Formulation :

Let H be a matrix of series describing a desired behavior, it is a reference model to achieve. Let CA^*B be the transfer of the system to be controlled, and P a filter such that $y = CA^*B \otimes P \otimes u$. The optimal filter such that $CA^*BPu \preceq Hu, \forall u$ is given by :

$$P_{opt} = (CA^*B) \oslash H$$

In manufacturing setting it is the one which delays as much as possible the input while ensuring that the output $y \preceq Hu \forall u$.

Optimal Filtering, model matching problem : JESA, Hardouin et al. 1995

Particular case : Neutral controller

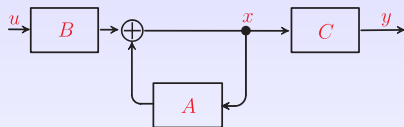
The model to match $H = CA^*B$, this means the output is not modified by the filter P but the input is delayed as much as possible.

Drawback :

it is also on open loop control, it doesn't take into account the possible disturbances acting on the system.

State Feedback controller synthesis

[Cottenceau 1999, Maia 2003, Lhommeau 2003]



system Equation :

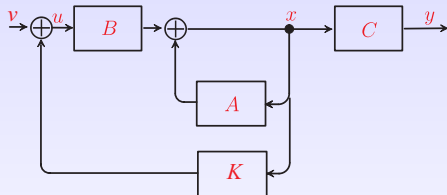
$$\begin{cases} \dot{x} = Ax \oplus Bu \\ y = Cx \end{cases}$$

input Output transfer function

$$y = CA^*Bu$$

State Feedback controller synthesis

[Cottenceau 1999, Maia 2003, Lhommeau 2003]



Controlled system :

$$\begin{cases} \dot{x} = Ax \oplus B(v \oplus Kx) \\ y = Cx \end{cases}$$

Closed loop system
transfer function :

$$y = C(A \oplus BK)^* Bv$$

Objective :

Compute the greatest
controller K such that :

$$C(A \oplus BK)^* B \preceq G_{ref}$$

Optimal State Feedback Controller :

► More

- Objective :

$$C(A \oplus BK)^* B \preceq G_{ref} \Leftrightarrow (A \oplus BK)^* B \preceq A^* B \oslash ((CA^* B) \oslash G_{ref})$$

- $\hat{G} = A^* B \oslash ((CA^* B) \oslash G_{ref})$ is the greatest transfer between u and x while preserving the constraint and belonging to $\text{Im}A^*B$ (i.e. which is achievable).

State Feedback controller synthesis

Proposition :

If $G_{ref} \succeq CA^*B = H$ then the greatest controller K ensuring $(A \oplus BK)\mathcal{G} \subset \mathcal{G}$ is given by :

$$\hat{K} = B \setminus \hat{\mathcal{G}} \phi \hat{\mathcal{G}}$$

Particular case : The neutral controller

If $G_{ref} = CA^*B$, the greatest controller is :

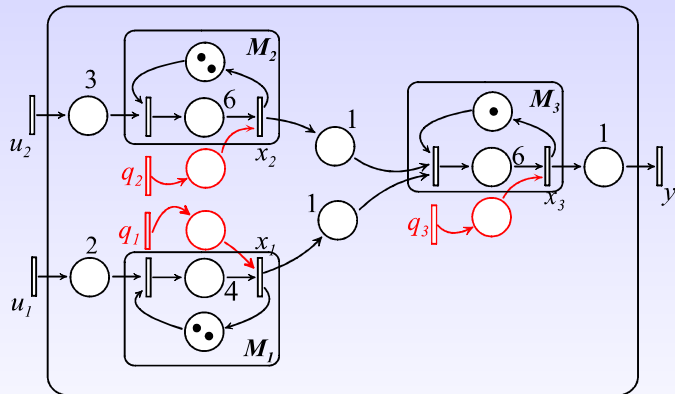
$$\hat{K} = (CA^*B) \setminus (CA^*B) \phi (A^*B)$$

it is the one which delays as much as possible the input while preserving the output.

Application : a kind of disturbances decoupling problem

Lagrange 2002, Lhommeau 2003

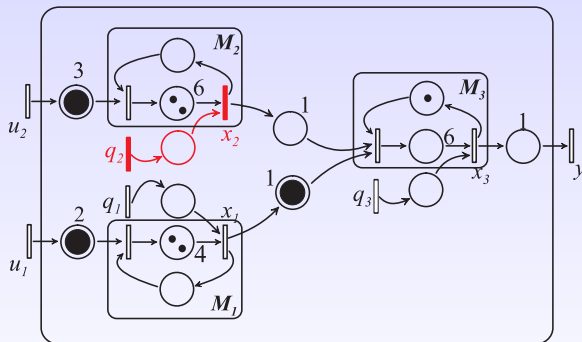
Problem statement :



Application : a kind of disturbances decoupling problem

Lagrange 2002, Lhommeau 2003

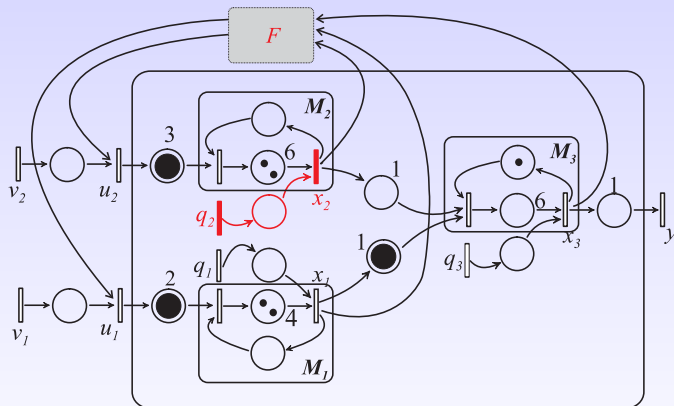
Problem statement :



Application : a kind of disturbances decoupling problem

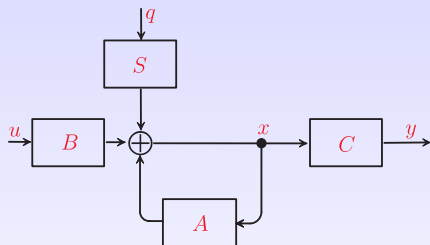
Lagrange 2002, Lhommeau 2003

Problem statement :



Application : a kind of disturbances decoupling problem

Lagrange 2002, Lhommeau 2003



State model :

$$x = Ax \oplus Bu \oplus Sq$$

$$y = Cx$$

Transfer Relation :

$$y = CA^*Bu \oplus CA^*Sq$$

Objective :

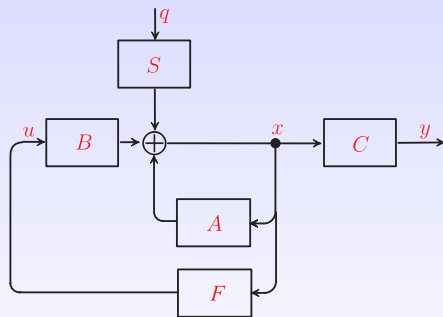
Compute the greatest controller $u = Fx$ such that

$$y = C(A \oplus BF)^*Sq$$

$$= CA^*Sq$$

Application : a kind of disturbances decoupling problem

Lagrange 2002, Lhommeau 2003



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Application : a kind of disturbances decoupling problem

Objective :

▸ definition

$$(A \oplus BF)^* S \stackrel{\ker C}{\equiv} A^* S$$

The greatest element of $[A^* S]_C$ [Cohen et al. 96] :

$$\Pi^C(A^* S) = C \setminus CA^* S$$

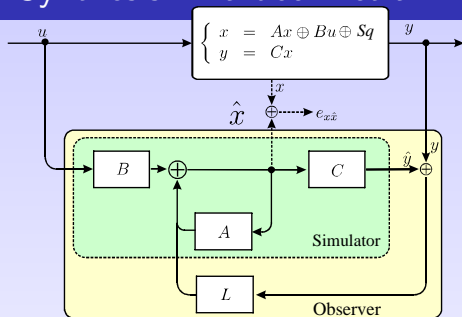
The greatest ideal principal A -invariant $(AV \subset V)$ included in $[A^* S]_C$:

$$\mathcal{G} = \{x \mid x \preceq \hat{\mathcal{G}} = A^* \setminus \Pi^C(A^* S)\}$$

The greatest controller F such that $(A \oplus BF)\mathcal{G} \subset \mathcal{G}$:

$$\hat{F} = B \setminus \hat{\mathcal{G}} \setminus \hat{\mathcal{G}}$$

Sub Observer Synthesis : Hardouin et al. IEEE TAC 2010

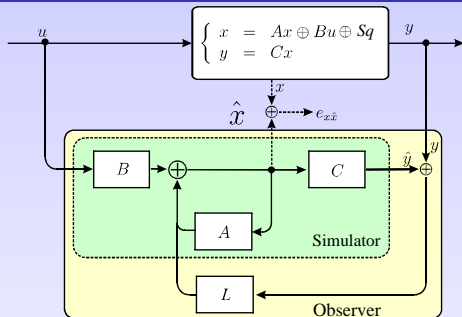


Objective :

Compute the greatest observer matrix L such that

$$\hat{x} \preceq x.$$

Sub Observer Synthesis :



System Equations :

► Matrix S

$$\dot{x} = Ax \oplus Bu \oplus Sq = A^*Bu \oplus A^*Sq$$

$$y = Cx = CA^*Bu \oplus CA^*Sq.$$

Estimated State Equations :

$$\dot{\hat{x}} = A\hat{x} \oplus Bu \oplus L(\hat{y} \oplus y)$$

$$\hat{y} = C\hat{x}.$$

Sub Observer Synthesis :

Constraints Satisfaction :

Compute the greatest observer matrix L such that

$$\begin{aligned}(A \oplus LC)^* Bu &\preceq A^* Bu && \forall u \\(A \oplus LC)^* LCA^* Sq &\preceq A^* Sq && \forall q,\end{aligned}$$

Constraints Satisfaction :

Compute the greatest matrix L such that

$$\begin{aligned}(A \oplus LC)^* B &\preceq A^* B \Leftrightarrow L \preceq (A^* B) \oslash (CA^* B) \\(A \oplus LC)^* LCA^* S &\preceq A^* S \Leftrightarrow L \preceq (A^* S) \oslash (CA^* S).\end{aligned}$$

Optimal Matrix :

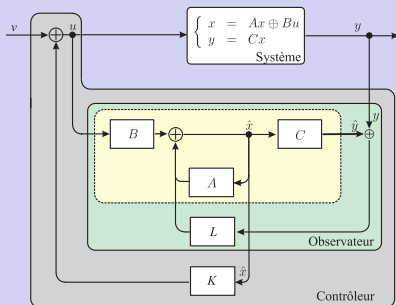
$$L_{opt} = ((A^*B) \oslash (CA^*B)) \wedge ((A^*S) \oslash (CA^*S))$$

is the greatest such that

$$\hat{x} \preceq x.$$

Control with an observer :

Principe :



Transfer Relation :

$$x = A^*B(K(A \oplus LC)^*B)^*v$$

$$y = CA^*B(K(A \oplus LC)^*B)^*v.$$

Control with an observer :

Objective :

$$\begin{aligned} CA^*B(K(A \oplus LC)^*B)^* &\preceq G_{ref} \\ \Leftrightarrow A^*B(K(A \oplus LC)^*B)^* &\preceq A^*B((CA^*B) \setminus G_{ref}) = \hat{\mathcal{G}} \end{aligned}$$

By considering the principal ideal : $\mathcal{G} = \{x | x \preceq \hat{\mathcal{G}}\}$

Controller \hat{K} :

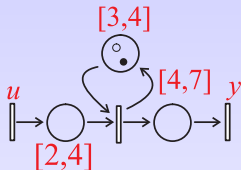
$$\hat{K} = B \setminus \hat{\mathcal{G}} \setminus \hat{\mathcal{G}} \setminus (A \oplus LC)^*$$

Properties :

Control strategy using both L_{opt} and \hat{K} is better (from the just in time point of view) than the one using only the output [Cottenceau et al. 99].

$(\max, +)$ -linear systems with uncertain parameters

[Lhommeau 2003]



This TEG can be represented as $I(\overline{\mathbb{Z}}_{\max}[\gamma])$

$$\begin{aligned}x(\gamma) &= [3\gamma^2, 4\gamma]x(\gamma) \oplus [2, 4]u(\gamma) \\y(\gamma) &= [4, 7]x(\gamma)\end{aligned}$$

The input/output transfer relation is given by :

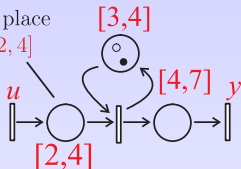
$$y(\gamma) = [6(3\gamma^2)^*, 11(4\gamma)^*]u(\gamma) = [\underline{H}, \overline{H}]u(\gamma) = \mathbf{H}u(\gamma)$$

All the transfer relations are included in interval $[\underline{H}, \overline{H}]$.

$(\max, +)$ -linear systems with uncertain parameters

[Lhommeau 2003]

temporisation associée à la place
comprise dans l'intervalle $[2, 4]$



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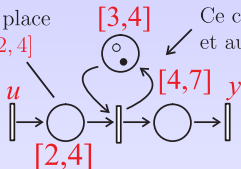
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$(\max, +)$ -linear systems with uncertain parameters

[Lhommeau 2003]

temporisation associée à la place
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Ce circuit contient au minimum 1 jeton
et au maximum 2 jetons.

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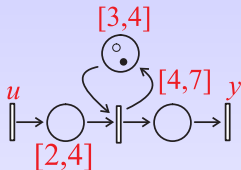
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$(\max, +)$ -linear systems with uncertain parameters

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The input/output transfer relation is given by :

$$y(\gamma) = [6(3\gamma^2)^*, 11(4\gamma)^*]u(\gamma) = [\underline{H}, \overline{H}]u(\gamma) = \mathbf{H}u(\gamma)$$

All the transfer relations are included in interval $[\underline{H}, \overline{H}]$.

Order relation in $I(\overline{\mathbb{Z}}_{\max}[\gamma])$

$$\mathbf{a} = [a, \bar{a}] \preceq_{I(\overline{\mathbb{Z}}_{\max}[\gamma])} \mathbf{b} = [b, \bar{b}] \Leftrightarrow a \preceq_{\overline{\mathbb{Z}}_{\max}[\gamma]} b \text{ et } \bar{a} \preceq_{\overline{\mathbb{Z}}_{\max}[\gamma]} \bar{b}$$

Residuation over $I(\overline{\mathbb{Z}}_{\max}[\gamma])$

The mapping $L_a : I(\overline{\mathbb{Z}}_{\max}[\gamma]) \rightarrow I(\overline{\mathbb{Z}}_{\max}[\gamma]), x \mapsto a \otimes x$ is residuated. The residual mapping L_a^\sharp is given by

$$L_a^\sharp(\mathbf{b}) = a \backslash \mathbf{b} = [a \backslash b \wedge \bar{a} \backslash \bar{b}, \bar{a} \backslash \bar{b}].$$

x is the greatest interval such that :

$$ax \preceq_{I(\overline{\mathbb{Z}}_{\max}[\gamma])} b$$

Order relation in $I(\overline{\mathbb{Z}}_{\max}[\gamma])$

$$\mathbf{a} = [a, \bar{a}] \preceq_{I(\overline{\mathbb{Z}}_{\max}[\gamma])} \mathbf{b} = [b, \bar{b}] \Leftrightarrow \underline{a} \preceq_{\overline{\mathbb{Z}}_{\max}[\gamma]} \underline{b} \text{ et } \bar{a} \preceq_{\overline{\mathbb{Z}}_{\max}[\gamma]} \bar{b}$$

Residuation over $I(\overline{\mathbb{Z}}_{\max}[\gamma])$

The mapping $\mathbf{L}_a : I(\overline{\mathbb{Z}}_{\max}[\gamma]) \rightarrow I(\overline{\mathbb{Z}}_{\max}[\gamma]), \mathbf{x} \mapsto \mathbf{a} \otimes \mathbf{x}$ is residuated. The residual mapping \mathbf{L}_a^\sharp is given by

$$\mathbf{L}_a^\sharp(\mathbf{b}) = \mathbf{a} \backslash \mathbf{b} = [\underline{a} \backslash \underline{b} \wedge \bar{a} \backslash \bar{b}, \bar{a} \backslash \bar{b}].$$

\mathbf{x} is the greatest interval such that :

$$\mathbf{a} \mathbf{x} \preceq_{I(\overline{\mathbb{Z}}_{\max}[\gamma])} \mathbf{b}$$

Conclusion :

Next works) :

- Application for High-Troughput Screening Systems
- How to deal with a model taking into account negative tokens, in order to take into account the transient behavior
- Control for system with maximal duration constraint, or dually a minimal event
- How to get a hierarchical approach, and also where put the sensors in an efficient way.
- and more

Scilab Toolboxes

- <http://www.istia.univ-angers.fr/~hardouin>
- <http://www.scilab.org/contrib/>
- <http://cermics.enpc.fr/~cohen-g//SED/book-online.html>

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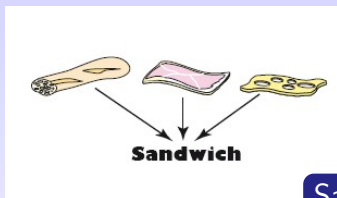
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Sandwiches Algebra [Cohen et al.]

1 piece of Bread + 1 slice of ham +
1 slice of cheese is equal to 1
sandwich. Another way of counting !

(max,+) algebra in few words

◀ Back

In classical algebra :

$$x = ax + b \Leftrightarrow x = (1 - a)^{-1}b$$

with the Taylor expansion (MacLaurin Series)

$$(1 - a)^{-1} \approx 1 + a + a^2 + a^3 + \dots$$

In semiring :

$x = ax \oplus b$ admits $x = a^*b$ as least solution. with $a^* = e \oplus a \oplus a^2 \oplus a^3 \oplus \dots$

$(\max, +)$ algebra in few words

Sketch of proof : if $a \in \text{Im } \mathcal{M}$ then $x^* \preceq a \Leftrightarrow x \preceq a^*$

◀ Back

$a \in \text{Im } \mathcal{M} \Rightarrow a = a^*$ hence $x^* \preceq a^* \Leftrightarrow (e \oplus x \oplus x^2 \oplus \dots) \preceq a^* \Leftrightarrow x \preceq a^*$

(max,+) algebra in few words

Sum of matrices $A \oplus B = C$

◀ Back

$$\begin{pmatrix} 2 & 5 \\ 3 & 7 \end{pmatrix} \oplus \begin{pmatrix} e & 8 \\ 1 & 3 \end{pmatrix} = \begin{pmatrix} 2 & 8 \\ 3 & 7 \end{pmatrix}$$

Product of matrices $A \otimes B = C$

◀ Back

$$\begin{pmatrix} 2 & 5 \\ \varepsilon & 3 \\ 1 & 8 \end{pmatrix} \otimes \begin{pmatrix} e \\ 1 \end{pmatrix} = \begin{pmatrix} 2 \otimes e \oplus 5 \otimes 1 \\ \varepsilon \otimes e \oplus 3 \otimes 1 \\ 1 \otimes e \oplus 8 \otimes 1 \end{pmatrix} = \begin{pmatrix} 6 \\ 4 \\ 9 \end{pmatrix}$$

(max,+) algebra in few words

Residuation of matrices $A \setminus B$ is the greatest solution of $A \otimes X \preceq B$

◀ Back

$$\begin{pmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{pmatrix} \otimes X \preceq \begin{pmatrix} 8 \\ 9 \\ 10 \end{pmatrix}$$

$$\begin{pmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{pmatrix} \setminus \begin{pmatrix} 8 \\ 9 \\ 10 \end{pmatrix} = \begin{pmatrix} (1 \setminus 8) \wedge (3 \setminus 9) \wedge (5 \setminus 10) \\ (2 \setminus 8) \wedge (4 \setminus 9) \wedge (6 \setminus 10) \end{pmatrix} = \begin{pmatrix} 5 \\ 4 \end{pmatrix}$$

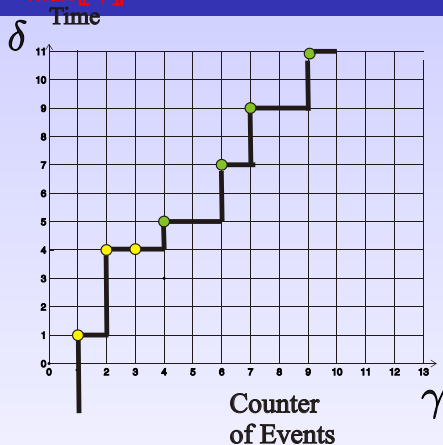
$\bar{\mathbb{Z}}_{\max}$

$\bar{\mathbb{Z}}_{\max} = (\mathbb{Z} \cup \{-\infty, +\infty\}, \max, +)$ is an idempotent semi ring.

$$(\forall a \in \bar{\mathbb{Z}}_{\max}, a \oplus a = a)$$

◀ Back

TEG Model in $\bar{\mathbb{Z}}_{\max}[\gamma]$

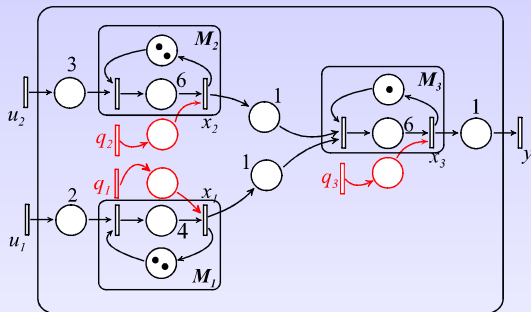


◀ Back

a series in $\bar{\mathbb{Z}}_{\max}[\gamma]$

$$s = \bigoplus_{k \in \mathbb{Z}} s(k) \gamma^k = 1\gamma \oplus 4\gamma^2 \oplus 5\gamma^4 \oplus 7\gamma^6 \oplus \dots$$

Sub Observer Synthesis :



Matrix S and input q :

► Back

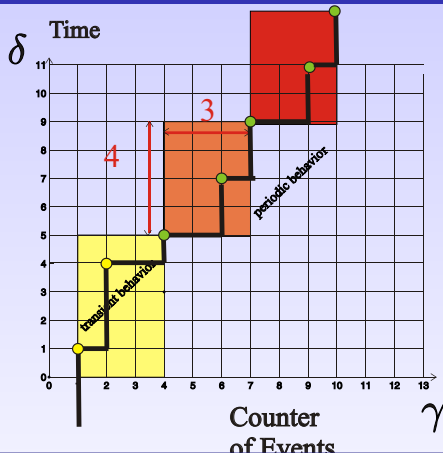
- vector q represents a vector of exogenous uncontrollable inputs (disturbance) which act on the system through matrix S .
- These disturbances lead to disable the transition firing, that is they decrease system performances and delay tokens output.

Computation $A^* = E \oplus A \oplus A^2 \oplus A^3 \oplus \dots$

$$A = \begin{pmatrix} e & \varepsilon \\ \varepsilon & e \end{pmatrix} \oplus \begin{pmatrix} \varepsilon & \gamma \\ 3 & \varepsilon \end{pmatrix} \oplus \begin{pmatrix} 3\gamma & \varepsilon \\ \varepsilon & 3\gamma \end{pmatrix} \oplus \begin{pmatrix} \varepsilon & 3\gamma^2 \\ 6\gamma & \varepsilon \end{pmatrix} \oplus \begin{pmatrix} 6\gamma^2 & \varepsilon \\ \varepsilon & 6\gamma^2 \end{pmatrix} \oplus \dots$$

Each entries is a periodic series $A^* = \begin{pmatrix} (3\gamma)^* & \gamma(3\gamma)^* \\ 3(3\gamma)^* & (3\gamma)^* \end{pmatrix}$

TEG Model in $\bar{Z}_{\max}[\gamma]$



◀ Back

a periodic series in $\bar{Z}_{\max}[\gamma]$

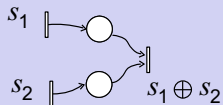
$$s = (1\gamma \oplus 4\gamma^2) \oplus (5\gamma^4 \oplus 7\gamma^6)(4\gamma^3)^*$$

The throughput is denoted by $\sigma_{\infty}(s) = 3/4$

Sum of series $s_1 \oplus s_2$

◀ Back

- Asymptotic slope $\sigma_\infty(s_1 \oplus s_2) = \min(\sigma_\infty(s_1), \sigma_\infty(s_2))$

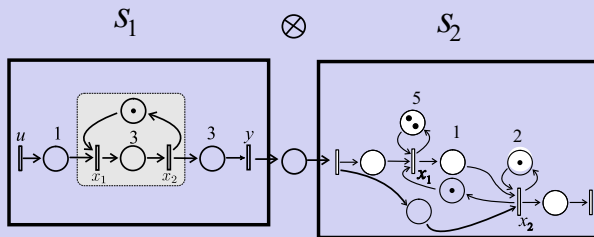


Cauchy product of series $s_1 \otimes s_2$

◀ Back

$$s_1 \otimes s_2 : (s_1 \otimes s_2)(k) = \bigoplus_{i+j=k} s_1(i) \otimes s_2(j).$$

- Asymptotic slope $\sigma_\infty(s_1 \otimes s_2) = \min(\sigma_\infty(s_1), \sigma_\infty(s_2))$



Details :

- Recall :

$$(a \oplus b)^* = a^*(ba^*)^*$$
$$a(ba)^* = (ab)^*a$$

- Let

$$C(A \oplus BK)^*B = CA^*(BKA^*)^*B = CA^*B(KA^*B)^*$$

.

- hence :

$$C(A \oplus BK)^*B \preceq G_{ref} \Leftrightarrow (KA^*B)^* \preceq ((CA^*B) \oslash G_{ref})$$
$$\Leftrightarrow (A \oplus BK)^*B \preceq A^*B((CA^*B) \oslash G_{ref}).$$

◀ Back

Application : a kind of disturbances decoupling problem

Definition [Davey et al.,1990],[Cohen et al, 1996]

▶ Back

The kernel of a mapping $C : \mathcal{X} \rightarrow \mathcal{Y}$, denoted $\ker C$, is defined by the following equivalence relation

$$x \stackrel{\ker C}{\equiv} x' \iff C(x) = C(x').$$